

The category of quantum graphs is closed

Bert Lindenhovius

Johannes Kepler University, Linz, Austria

July 6, 2026



Joint work with Andre Kornell
based on the preprint

Quantum graphs of homomorphisms

Motivation

Two notions of quantum graph homomorphisms have developed independently:

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Are these manifestations of the same mathematical structure?

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Are these manifestations of the same mathematical structure?

Goals:

- Define a category of quantum graphs;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Are these manifestations of the same mathematical structure?

Goals:

- Define a category of quantum graphs;
- Describe both meanings of quantum graph homomorphisms in this category;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Are these manifestations of the same mathematical structure?

Goals:

- Define a category of quantum graphs;
- Describe both meanings of quantum graph homomorphisms in this category;
- Show that the category is monoidal closed;

Motivation

Two notions of quantum graph homomorphisms have developed independently:

- Homomorphisms of quantum graphs;
- Quantum homomorphisms of ordinary graphs;

Are these manifestations of the same mathematical structure?

Goals:

- Define a category of quantum graphs;
- Describe both meanings of quantum graph homomorphisms in this category;
- Show that the category is monoidal closed;
- Use the monoidal closure to characterize winning quantum strategies of graph homomorphism games.

Classical graphs

Definition

- A graph is a pair $G = (V_G, e_G)$ consisting of:
 - ▶ a set V_G of vertices;
 - ▶ a symmetric relation e_G on V_G of edges (write $g_1 \sim g_2 \iff (g_1, g_2) \in e_G$);
- A homomorphism $G \rightarrow H$ is a map $\varphi : V_G \rightarrow V_H$ such that

$$g_1 \sim g_2 \implies \varphi(g_1) \sim \varphi(g_2).$$

We denote the category of graphs and homomorphisms by **Gph**.

Classical graphs

Definition

- A graph is a pair $G = (V_G, e_G)$ consisting of:
 - ▶ a set V_G of vertices;
 - ▶ a symmetric relation e_G on V_G of edges (write $g_1 \sim g_2 \iff (g_1, g_2) \in e_G$);
- A homomorphism $G \rightarrow H$ is a map $\varphi : V_G \rightarrow V_H$ such that

$$g_1 \sim g_2 \implies \varphi(g_1) \sim \varphi(g_2).$$

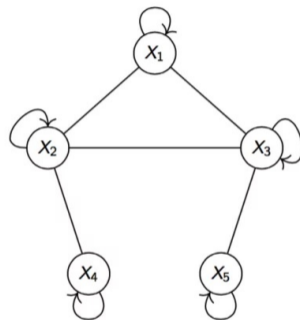
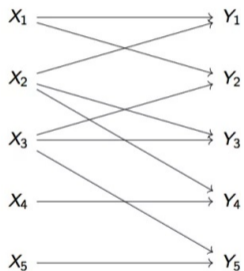
We denote the category of graphs and homomorphisms by **Gph**.

Remarks

- Graphs are allowed to have loops, but are possibly loopless.
- G has loops on every vertex $\iff e_G$ is reflexive.
- G has no loops at all $\iff e_G$ is irreflexive.

Confusibility graphs

Input messages (X) $\xrightarrow{\Phi}$ Output messages (Y)



Quantum confusibility graphs

Definition (Duan, Severini, Winter)

Let $\varphi : M_m(\mathbb{C}) \rightarrow M_n(\mathbb{C})$ be a quantum channel (a ctp map) with Kraus matrices $(K_i)_{i \in I}$. Then the quantum confusibility graph of φ is the subspace

$$\mathcal{G}_\varphi = \text{span}\{K_i^\dagger K_j : i, j \in I\} \subseteq M_m(\mathbb{C}).$$

Quantum confusibility graphs

Definition (Duan, Severini, Winter)

Let $\varphi : M_m(\mathbb{C}) \rightarrow M_n(\mathbb{C})$ be a quantum channel (a cptp map) with Kraus matrices $(K_i)_{i \in I}$. Then the quantum confusibility graph of φ is the subspace

$$\mathcal{G}_\varphi = \text{span}\{K_i^\dagger K_j : i, j \in I\} \subseteq M_m(\mathbb{C}).$$

Observations

- \mathcal{G}_φ is independent of the choice of Kraus matrices;
- For pure states p_x and p_y (for unit vectors $x, y \in \mathbb{C}^m$), the following are equivalent:
 - ▶ φ confuses inputs p_x and p_y ;
 - ▶ $\text{tr}(\varphi(p_x)\varphi(p_y)) \neq 0$;
 - ▶ $x^\dagger(\mathcal{G}_\varphi)y \neq \{0\}$.
- \mathcal{G}_φ is selfadjoint and unital.

Quantum confusibility graphs

Definition (Duan, Severini, Winter)

Let $\varphi : M_m(\mathbb{C}) \rightarrow M_n(\mathbb{C})$ be a quantum channel (a cptp map) with Kraus matrices $(K_i)_{i \in I}$. Then the quantum confusibility graph of φ is the subspace

$$\mathcal{G}_\varphi = \text{span}\{K_i^\dagger K_j : i, j \in I\} \subseteq M_m(\mathbb{C}).$$

Observations

- \mathcal{G}_φ is independent of the choice of Kraus matrices;
- For pure states p_x and p_y (for unit vectors $x, y \in \mathbb{C}^m$), the following are equivalent:
 - ▶ φ confuses inputs p_x and p_y ;
 - ▶ $\text{tr}(\varphi(p_x)\varphi(p_y)) \neq 0$;
 - ▶ $x^\dagger(\mathcal{G}_\varphi)y \neq \{0\}$.
- \mathcal{G}_φ is selfadjoint and unital.

Definition

A finite (reflexive) quantum graph is a (unital) selfadjoint subspace of $M_m(\mathbb{C})$.

Graph homomorphism games

Definition

Let G and H be finite graphs. The (G, H) -homomorphism game is played by:

- 1 A verifier who selects a pair of vertices $g_1, g_2 \in V_G$;
- 2 Alice and Bob, who selects a vertex $h_1 \in V_H$ and a vertex $h_2 \in V_H$, respectively;
- 3 the winning condition for Alice and Bob is $(g_1 = g_2 \Rightarrow h_1 = h_2) \ \& \ (g_1 \sim g_2 \Rightarrow h_1 \sim h_2)$.

Graph homomorphism games

Definition

Let G and H be finite graphs. The (G, H) -homomorphism game is played by:

- 1 A verifier who selects a pair of vertices $g_1, g_2 \in V_G$;
- 2 Alice and Bob, who selects a vertex $h_1 \in V_H$ and a vertex $h_2 \in V_H$, respectively;
- 3 the winning condition for Alice and Bob is $(g_1 = g_2 \Rightarrow h_1 = h_2) \ \& \ (g_1 \sim g_2 \Rightarrow h_1 \sim h_2)$.

Classical strategy (informal)

Alice and Bob agree in advance, possibly using shared randomness, on a single function $V_G \rightarrow V_H$, and each applies it locally without communication.

Graph homomorphism games

Definition

Let G and H be finite graphs. The (G, H) -homomorphism game is played by:

- 1 A verifier who selects a pair of vertices $g_1, g_2 \in V_G$;
- 2 Alice and Bob, who selects a vertex $h_1 \in V_H$ and a vertex $h_2 \in V_H$, respectively;
- 3 the winning condition for Alice and Bob is $(g_1 = g_2 \Rightarrow h_1 = h_2) \ \& \ (g_1 \sim g_2 \Rightarrow h_1 \sim h_2)$.

Classical strategy (informal)

Alice and Bob agree in advance, possibly using shared randomness, on a single function $V_G \rightarrow V_H$, and each applies it locally without communication.

Theorem

For finite simple graphs G and H the following are equivalent:

- *There is a classical winning strategy for the (G, H) -homomorphism game;*
- *There is a homomorphism $G \rightarrow H$.*

Monoidal structures on **Gph**

- Direct product $G \times H$:

Monoidal structures on **Gph**

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;

Monoidal structures on **Gph**

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;

Monoidal structures on **Gph**

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
 - ▶ Monoidal unit: $\overline{K_1}$ (one vertex, one loop);

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
 - ▶ Monoidal unit: $\overline{K_1}$ (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:

Monoidal structures on **Gph**

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
 - ▶ Monoidal unit: $\overline{K_1}$ (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,

Monoidal structures on **Gph**

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
 - ▶ Monoidal unit: $\overline{K_1}$ (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2$ and $h_1 \sim h_2$;
 - ▶ Monoidal unit: $\overline{K_1}$ (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.
- Box product $G \square H$:

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2$ and $h_1 \sim h_2$;
 - ▶ Monoidal unit: K_1 (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.
- Box product $G \square H$:
 - ▶ $V_{G \square H} = V_G \times V_H$;

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
- ▶ Monoidal unit: K_1 (one vertex, no loop).

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
- ▶ Monoidal unit: K_1 (one vertex, no loop).
- ▶ Internal hom $[G, H]_{\square}$:

Monoidal structures on **Gph**

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
- ▶ Monoidal unit: K_1 (one vertex, no loop).
- ▶ Internal hom $[G, H]_{\square}$:
 - ★ $V_{[G, H]_{\square}} = \mathbf{Gph}(G, H)$;

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
- ▶ Monoidal unit: K_1 (one vertex, no loop).
- ▶ Internal hom $[G, H]_{\square}$:
 - ★ $V_{[G, H]_{\square}} = \mathbf{Gph}(G, H)$;
 - ★ $\varphi_1 \sim_{[G, H]_{\square}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:

- ▶ $V_{G \times H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
- ▶ Monoidal unit: K_1 (one vertex, one loop);
- ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

- Box product $G \square H$:

- ▶ $V_{G \square H} = V_G \times V_H$;
- ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
- ▶ Monoidal unit: K_1 (one vertex, no loop).
- ▶ Internal hom $[G, H]_{\square}$:
 - ★ $V_{[G, H]_{\square}} = \mathbf{Gph}(G, H)$;
 - ★ $\varphi_1 \sim_{[G, H]_{\square}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

Monoidal structures on \mathbf{Gph}

- Direct product $G \times H$:
 - ▶ $V_{G \times H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \times H} (g_2, h_2) \iff g_1 \sim g_2 \text{ and } h_1 \sim h_2$;
 - ▶ Monoidal unit: K_1 (one vertex, one loop);
 - ▶ Internal hom $[G, H]_{\times}$:
 - ★ $V_{[G, H]_{\times}} = (V_H)^{V_G}$,
 - ★ $\varphi_1 \sim_{[G, H]_{\times}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.
- Box product $G \square H$:
 - ▶ $V_{G \square H} = V_G \times V_H$;
 - ▶ $(g_1, h_1) \sim_{G \square H} (g_2, h_2)$ iff $(g_1 = g_2 \text{ and } h_1 \sim h_2)$ or $(g_1 \sim g_2 \text{ and } h_1 = h_2)$;
 - ▶ Monoidal unit: K_1 (one vertex, no loop).
 - ▶ Internal hom $[G, H]_{\square}$:
 - ★ $V_{[G, H]_{\square}} = \mathbf{Gph}(G, H)$;
 - ★ $\varphi_1 \sim_{[G, H]_{\square}} \varphi_2$ iff $\varphi_1(g) \sim \varphi_2(g)$ for all $g \in V_G$.

Theorem

A homomorphism $\varphi : G \rightarrow H$ exists $\iff [G, H]_{\square} \neq \emptyset$.

Quantum strategies

Definition (informal)

A quantum strategy is:

- Alice and Bob share an entangled quantum state in advance,
- Alice measures it with outcome h_1 , choosing her measurement based on g_1 ;
- Bob measures it with outcome h_2 , choosing his measurement based on g_2 .

Quantum strategies

Definition (informal)

A quantum strategy is:

- Alice and Bob share an entangled quantum state in advance,
- Alice measures it with outcome h_1 , choosing her measurement based on g_1 ;
- Bob measures it with outcome h_2 , choosing his measurement based on g_2 .

Winning quantum strategies

- Found by Mančinska, Roberson (2016), called 'quantum homomorphism';
- Can exist even when there is no winning classical strategy.
- Quantum homomorphisms formalized by Musto, Reutter, and Verdon (2018).

Quantum sets

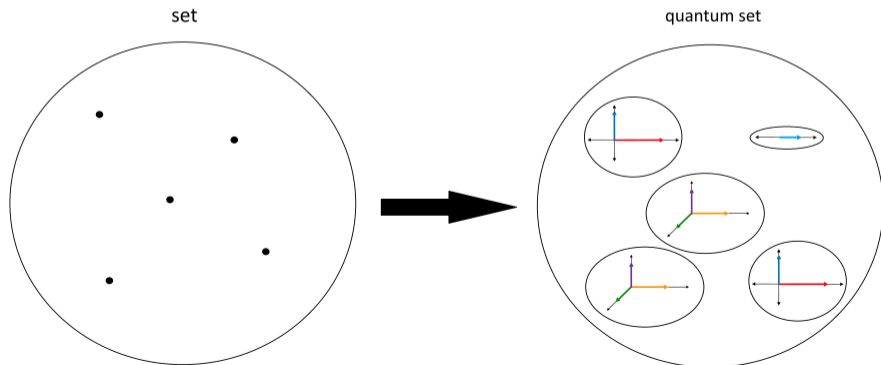
Definition

- A quantum set \mathcal{X} is a collection of Hilbert spaces (called atoms) of finite dimension > 0 ;
- $\ell^\infty(\mathcal{X}) := \bigoplus_{X \in \mathcal{X}}^\infty B(X)$ is its associated hereditarily atomic von Neumann algebra;

Quantum sets

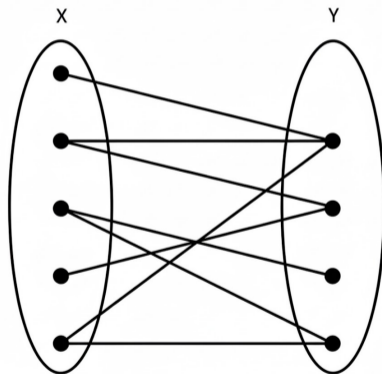
Definition

- A quantum set \mathcal{X} is a collection of Hilbert spaces (called atoms) of finite dimension > 0 ;
- $\ell^\infty(\mathcal{X}) := \bigoplus_{X \in \mathcal{X}}^\infty B(X)$ is its associated hereditarily atomic von Neumann algebra;



Binary relations

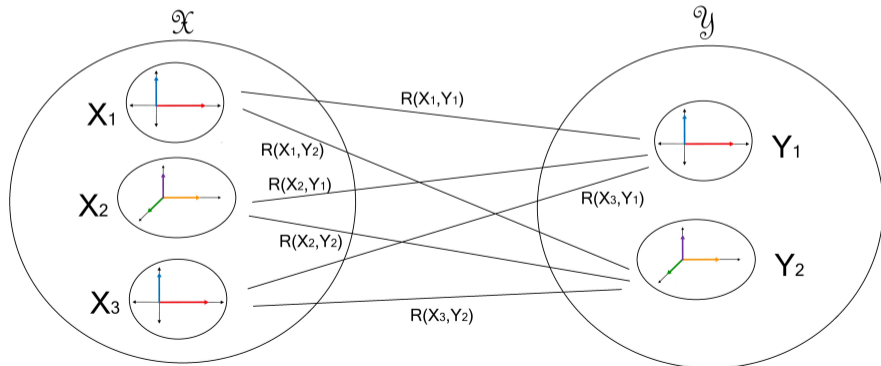
A binary relation $r : X \rightarrow Y$ is a bipartite graph (possibly infinite vertices).



Relations between quantum sets

Definition

A relation $R : \mathcal{X} \rightarrow \mathcal{Y}$ between quantum sets is a complete bipartite graph, where the edge between atoms $X \in \mathcal{X}$ and $Y \in \mathcal{Y}$ is labeled by a space $R(X, Y)$ of linear operators $X \rightarrow Y$. We denote the category of quantum sets and relations by **qRel**.



Connections with ordinary sets and von Neumann algebras

- Fully faithful functor $\ell^\infty : \mathbf{qRel} \rightarrow \mathbf{WRel}$, the category of von Neumann algebras and Weaver's quantum relations;
- Fully faithful functor $'(-) : \mathbf{Rel} \rightarrow \mathbf{qRel}$, $S \mapsto (\mathbb{C} : s \in S)$.

Discrete quantization

- **Quantization**: finding noncommutative generalizations of mathematical structures;

Discrete quantization

- **Quantization**: finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**

Discrete quantization

- **Quantization**: finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.
- **Mathematical structures:** sets and relations satisfying constraints in terms of \leq and $(-)^\dagger$.

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.
- **Mathematical structures:** sets and relations satisfying constraints in terms of \leq and $(-)^{\dagger}$.
- **Discrete quantization:**

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.
- **Mathematical structures:** sets and relations satisfying constraints in terms of \leq and $(-)^{\dagger}$.
- **Discrete quantization:**
 - ▶ sets \rightarrow quantum sets;

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.
- **Mathematical structures:** sets and relations satisfying constraints in terms of \leq and $(-)^{\dagger}$.
- **Discrete quantization:**
 - ▶ sets \rightarrow quantum sets;
 - ▶ relations between sets \rightarrow relations between quantum sets;

Discrete quantization

- **Quantization:** finding noncommutative generalizations of mathematical structures;
- **Shared properties of homsets of Rel and qRel:**
 - ▶ Reversing relations $R : \mathcal{X} \rightarrow \mathcal{Y} \implies R^\dagger : \mathcal{Y} \rightarrow \mathcal{X}$;
 - ▶ Ordering relations $R \leq S : \mathcal{X} \rightarrow \mathcal{Y}$.
- **Mathematical structures:** sets and relations satisfying constraints in terms of \leq and $(-)^{\dagger}$.
- **Discrete quantization:**
 - ▶ sets \rightarrow quantum sets;
 - ▶ relations between sets \rightarrow relations between quantum sets;
 - ▶ same constraints.

Example: functions between quantum sets

Observation

A function $f : X \rightarrow Y$ is a morphism in **Rel** such that

$$f^\dagger \circ f \geq \text{id}_X, \quad f \circ f^\dagger \leq \text{id}_Y.$$

Example: functions between quantum sets

Observation

A function $f : X \rightarrow Y$ is a morphism in **Rel** such that

$$f^\dagger \circ f \geq \text{id}_X, \quad f \circ f^\dagger \leq \text{id}_Y.$$

Definition

A function $F : \mathcal{X} \rightarrow \mathcal{Y}$ between quantum sets is a relation such that

$$F^\dagger \circ F \geq I_{\mathcal{X}}, \quad F \circ F^\dagger \leq I_{\mathcal{Y}}.$$

We denote the category of quantum sets and functions by **qSet**.

Properties of **qSet**

Theorem (Kornell (2020))

- $\ell^\infty : \mathbf{qRel} \rightarrow \mathbf{WRel} \xrightarrow{\text{restricts and corestricts}} \text{fully faithful } \ell^\infty : \mathbf{qSet} \rightarrow \mathbf{WStar}^{\text{op}}$,
 $\mathbf{WStar} = \text{von Neumann algebras and normal unital } *- \text{homomorphisms.}$
- $'(-) : \mathbf{Rel} \rightarrow \mathbf{qRel} \xrightarrow{\text{restricts and corestricts}} \text{fully faithful } '(-) : \mathbf{Set} \rightarrow \mathbf{qSet}.$
- \mathbf{qSet} is generated by the quantum sets $\mathcal{Q}_n = \{\mathbb{C}^n\}$.
- \mathbf{qSet} is symmetric monoidal closed:
 - ▶ product $\mathcal{X} \times \mathcal{Y} = \{X \otimes Y : X \in \mathcal{X}, Y \in \mathcal{Y}\}$ is semicartesian;
 - ▶ Monoidal unit: \mathcal{Q}_1 ;
 - ▶ Monoidal closure $\mathcal{Y}^{\mathcal{X}}$ constructed in terms of representations.

Discrete quantization is a unifying framework:

- Quantum graphs^{1 2} → quantum error correction, nonlocal games;
- Quantum posets^{3 4}
 - ▶ Quantum cpos⁴ → recursive pure quantum computing;
 - ▶ Quantum suplattices → quantum model of full linear logic?
→ quantum model of differential logic?
- Quantum metric spaces^{5 6} → quantum error correction;

¹Duan, Severini, Winter (2013)

²Mančinska, Roberson (2016)

³Weaver (2010)

⁴Kornell, L., Mislove (2026)

⁵Kuperberg, Weaver (2010)

⁶Weaver (2017)

Quantum graphs

Definition of the category **Gph** of graphs and homomorphisms in relational terms:

- Objects: pairs $G = (V_G, e_G)$ with a set V_G and a relation e_G on V_G such that $e_G^\dagger = e_G$.
- Homomorphisms $\varphi : G \rightarrow H$ are maps $V_G \rightarrow V_H$ such that $\varphi \circ e_G \leq e_H \circ \varphi$.

Quantum graphs

Definition of the category **Gph** of graphs and homomorphisms in relational terms:

- Objects: pairs $G = (V_G, e_G)$ with a set V_G and a relation e_G on V_G such that $e_G^\dagger = e_G$.
- Homomorphisms $\varphi : G \rightarrow H$ are maps $V_G \rightarrow V_H$ such that $\varphi \circ e_G \leq e_H \circ \varphi$.

Definition of the category **qGph** of quantum graphs and homomorphisms:

- Objects: $G = (\mathcal{V}_G, E_G)$, where \mathcal{V}_G is a quantum set and E_G a relation on \mathcal{V}_G such that $E_G^\dagger = E_G$.
- Homomorphisms $\Phi : G \rightarrow H$ are maps $\mathcal{V}_G \rightarrow \mathcal{V}_H$ such that $\Phi \circ E_G \leq E_H \circ \Phi$.

Quantum graphs

Definition of the category **Gph** of graphs and homomorphisms in relational terms:

- Objects: pairs $G = (V_G, e_G)$ with a set V_G and a relation e_G on V_G such that $e_G^\dagger = e_G$.
- Homomorphisms $\varphi : G \rightarrow H$ are maps $V_G \rightarrow V_H$ such that $\varphi \circ e_G \leq e_H \circ \varphi$.

Definition of the category **qGph** of quantum graphs and homomorphisms:

- Objects: $G = (\mathcal{V}_G, E_G)$, where \mathcal{V}_G is a quantum set and E_G a relation on \mathcal{V}_G such that $E_G^\dagger = E_G$.
- Homomorphisms $\Phi : G \rightarrow H$ are maps $\mathcal{V}_G \rightarrow \mathcal{V}_H$ such that $\Phi \circ E_G \leq E_H \circ \Phi$.

Example (Quantum confusibility graphs)

Let $\varphi : M_m(\mathbb{C}) \rightarrow M_n(\mathbb{C})$ be a quantum channel. Then $G_\varphi = (\mathcal{Q}_m, E_\varphi)$ with $E_\varphi(\mathbb{C}^m, \mathbb{C}^m) = \mathcal{G}_\varphi$.

Quantum graphs

Definition of the category **Gph** of graphs and homomorphisms in relational terms:

- Objects: pairs $G = (V_G, e_G)$ with a set V_G and a relation e_G on V_G such that $e_G^\dagger = e_G$.
- Homomorphisms $\varphi : G \rightarrow H$ are maps $V_G \rightarrow V_H$ such that $\varphi \circ e_G \leq e_H \circ \varphi$.

Definition of the category **qGph** of quantum graphs and homomorphisms:

- Objects: $G = (\mathcal{V}_G, E_G)$, where \mathcal{V}_G is a quantum set and E_G a relation on \mathcal{V}_G such that $E_G^\dagger = E_G$.
- Homomorphisms $\Phi : G \rightarrow H$ are maps $\mathcal{V}_G \rightarrow \mathcal{V}_H$ such that $\Phi \circ E_G \leq E_H \circ \Phi$.

Example (Quantum confusibility graphs)

Let $\varphi : M_m(\mathbb{C}) \rightarrow M_n(\mathbb{C})$ be a quantum channel. Then $G_\varphi = (\mathcal{Q}_m, E_\varphi)$ with $E_\varphi(\mathbb{C}^m, \mathbb{C}^m) = \mathcal{G}_\varphi$.

Theorem

There is a fully faithful functor $'(-) : \mathbf{Gph} \rightarrow \mathbf{qGph}$, $(V, e) \mapsto ('V, 'e)$.

Quantum functions

Observation

There is a bijection between:

- Families of functions $X \rightarrow Y$ indexed by a set Z ;
- Functions $X \times Z \rightarrow Y$.

Quantum functions

Observation

There is a bijection between:

- Families of functions $X \rightarrow Y$ indexed by a set Z ;
- Functions $X \times Z \rightarrow Y$.

Definition

A quantum family of functions $\mathcal{X} \rightarrow \mathcal{Y}$ indexed by a quantum set \mathcal{Z} is a function $\mathcal{X} \times \mathcal{Z} \rightarrow \mathcal{Y}$.

Quantum functions

Observation

There is a bijection between:

- Families of functions $X \rightarrow Y$ indexed by a set Z ;
- Functions $X \times Z \rightarrow Y$.

Definition

A quantum family of functions $\mathcal{X} \rightarrow \mathcal{Y}$ indexed by a quantum set \mathcal{Z} is a function $\mathcal{X} \times \mathcal{Z} \rightarrow \mathcal{Y}$.

Definition

An n -dimensional quantum function $\mathcal{X} \rightarrow \mathcal{Y}$ is a quantum family of functions $\mathcal{X} \rightarrow \mathcal{Y}$ indexed by \mathcal{Q}_n .

Observations

- For any quantum set \mathcal{X} there is a bijection between:

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:
 - ▶ functions $A \rightarrow B$;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:
 - ▶ functions $A \rightarrow B$;
 - ▶ one-dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:
 - ▶ functions $A \rightarrow B$;
 - ▶ one-dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$;
- So we have an identification between:

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:
 - ▶ functions $A \rightarrow B$;
 - ▶ one-dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$;
- So we have an identification between:
 - ▶ the function set B^A ;

Observations

- For any quantum set \mathcal{X} there is a bijection between:
 - ▶ functions $\mathcal{Q}_1 \rightarrow \mathcal{X}$;
 - ▶ one-dimensional atoms of \mathcal{X} (the classical part of \mathcal{X});
- More generally, there is a bijection between:
 - ▶ (equivalence classes of) functions $\mathcal{Q}_n \rightarrow \mathcal{X}$;
 - ▶ n -dimensional atoms of \mathcal{X} ;
- For any quantum sets \mathcal{X} and \mathcal{Y} there is a bijection between
 - ▶ n -dimensional quantum functions $\mathcal{X} \rightarrow \mathcal{Y}$;
 - ▶ n -dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$.
- For $\mathcal{X} = 'A$ and $\mathcal{Y} = 'B$, we have a bijection between:
 - ▶ functions $A \rightarrow B$;
 - ▶ one-dimensional atoms of $\mathcal{Y}^{\mathcal{X}}$;
- So we have an identification between:
 - ▶ the function set B^A ;
 - ▶ the classical part of the quantum function set $('B)^{'A}$.

Monoidal products on **Gph** in relational terms

Direct product $G \times H$:

- $V_{G \times H} = V_G \times V_H$;
- $e_{G \times H} = e_G \times e_H$;
- Monoidal unit: \overline{K}_1 (one vertex, one loop);

Box product $G \square H$:

- $V_{G \square H} = V_G \times V_H$;
- $e_{G \square H} = (e_G \times \text{id}_H) \vee (\text{id}_G \times e_H)$;
- Monoidal unit: K_1 (one vertex, no loops).

Monoidal products on \mathbf{qGph}

Direct product $G \times H$:

- $\mathcal{V}_{G \times H} = \mathcal{V}_G \times \mathcal{V}_H$;
- $E_{G \times H} = E_G \times E_H$;
- Monoidal unit: $\bar{Q}_1 := (Q_1, I_{Q_1}) = 'K_1$ (one vertex, one loop);

Box product $G \square H$:

- $\mathcal{V}_{G \square H} = \mathcal{V}_G \times \mathcal{V}_H$;
- $E_{G \square H} = (E_G \times I_H) \vee (I_G \times E_H)$;
- Monoidal unit: $Q_1 := (Q_1, 0) = 'K_1$ (one vertex, no loops);

Quantum homomorphisms

Let $Q_n = (\mathcal{Q}_n, 0)$ and $\overline{Q}_n = (\mathcal{Q}_n, I_{\mathcal{Q}_n})$.

Definition

An n -dimensional quantum homomorphism $G \rightarrow H$ between (quantum) graphs is an n -dimensional quantum function $\Phi : \mathcal{V}_G \rightarrow \mathcal{V}_H$ that satisfies the following equivalent conditions:

- $\Phi : G \times \overline{Q}_n \rightarrow H$ is a homomorphism;
- $\Phi : G \square Q_n \rightarrow H$ is a homomorphism.

Quantum homomorphisms

Let $Q_n = (Q_n, 0)$ and $\overline{Q}_n = (Q_n, I_{Q_n})$.

Definition

An n -dimensional quantum homomorphism $G \rightarrow H$ between (quantum) graphs is an n -dimensional quantum function $\Phi : \mathcal{V}_G \rightarrow \mathcal{V}_H$ that satisfies the following equivalent conditions:

- $\Phi : G \times \overline{Q}_n \rightarrow H$ is a homomorphism;
- $\Phi : G \square Q_n \rightarrow H$ is a homomorphism.

Observations

- These are precisely the quantum homomorphisms of Musto, Reutter, and Verdon (2018);
- Hence, these quantum homomorphisms correspond to winning quantum strategies;
- 1-dimensional quantum homomorphisms = homomorphisms.

The internal hom of \mathbf{qGph}

Theorem (Kornell, L.)

$(\mathbf{qGph}, \square, Q_1)$ is closed.

The internal hom of \mathbf{qGph}

Theorem (Kornell, L.)

$(\mathbf{qGph}, \square, Q_1)$ is closed.

Proof (sketch).

Exploit that \mathbf{qSet} is closed.

- 1 $\mathcal{V}_{[G,H]}$ is the largest subset \mathcal{W} of the quantum function set $(\mathcal{V}_H)^{\mathcal{V}_G}$ such that $\text{Eval}_{G,H} \circ (\text{Inc}_{\mathcal{W}} \times I_G)$ is a homomorphism $(\mathcal{W} \times \mathcal{V}_G, I_{\mathcal{W}} \times E_G) \rightarrow H$.
- 2 $E_{[G,H]}$ is the largest symmetric relation R on $\mathcal{V}_{[G,H]} = \mathcal{W}$ such that $\text{Eval}_{G,H} \circ (\text{Inc}_{\mathcal{W}} \times I_G)$ is a homomorphism $(\mathcal{W} \times \mathcal{V}_G, R \times I_G) \rightarrow H$.

The internal hom of \mathbf{qGph}

Theorem (Kornell, L.)

$(\mathbf{qGph}, \square, Q_1)$ is closed.

Proof (sketch).

Exploit that \mathbf{qSet} is closed.

- 1 $\mathcal{V}_{[G,H]}$ is the largest subset \mathcal{W} of the quantum function set $(\mathcal{V}_H)^{\mathcal{V}_G}$ such that $\text{Eval}_{G,H} \circ (\text{Inc}_{\mathcal{W}} \times I_G)$ is a homomorphism $(\mathcal{W} \times \mathcal{V}_G, I_{\mathcal{W}} \times E_G) \rightarrow H$.
- 2 $E_{[G,H]}$ is the largest symmetric relation R on $\mathcal{V}_{[G,H]} = \mathcal{W}$ such that $\text{Eval}_{G,H} \circ (\text{Inc}_{\mathcal{W}} \times I_G)$ is a homomorphism $(\mathcal{W} \times \mathcal{V}_G, R \times I_G) \rightarrow H$.

Observations

- Proof for $(\mathbf{Gph}, \square, K_1)$ exactly the same.
- We have a bijection between:
 - ▶ n -dimensional quantum homomorphisms $G \rightarrow H$;
 - ▶ n -dimensional atoms of $[G, H]$;

Winning quantum strategies

Theorem (Kornell, L.)

For finite simple graphs G and H the following are equivalent:

- *There is a winning quantum strategy for the (G, H) -homomorphism game;*
- *There is a quantum homomorphism $G \rightarrow H$;*
- *$[G, H] \neq \emptyset$.*

Winning quantum strategies

Theorem (Kornell, L.)

For finite simple graphs G and H the following are equivalent:












- There is a winning quantum strategy for the (G, H) -homomorphism game;
- There is a quantum homomorphism $G \rightarrow H$;
- $[G, H] \neq \emptyset$.

Corollary

The following are equivalent for graphs G and H :

- There is a quantum homomorphism $G \rightarrow H$, but no homomorphism $G \rightarrow H$.
- $[G, H] \neq \emptyset$, but there are no one-dimensional atoms.

References

-  R. Duan, S. Severini, A. Winter, *Zero-error communication via quantum channels, noncommutative graphs, and a quantum Lovász number*. IEEE Trans. Inf. Theory, 2013.
-  G. Jenča, B.L., *Quantum Suplattices*, Proceedings of QPL, 2023.
-  A. Kornell, *Quantum Sets*, Journal of Mathematical Physics, 2020.
-  A. Kornell, B.L., *Quantum graphs of homomorphisms*, preprint (2026).
-  A. Kornell, B.L., M. Mislove, *A Category of Quantum Posets*, Indagationes Mathematicae, 2022.
-  A. Kornell, B.L., M. Mislove, *Categories of Quantum CPOs*, MSCS, 2026.
-  G. Kuperberg, N. Weaver, *A von Neumann algebra approach to quantum metrics*, AMS, 2012.
-  L. Mančinska, D.E. Roberson, *Quantum homomorphisms*, J. Comb. Theory, 2016.
-  B. Musto, D. Reutter, D. Verdon, *A compositional approach to quantum functions*. J. Math. Phys., 2018.
-  N. Weaver, *Quantum graphs as quantum relations*, J. Geom. Anal., 2021.
-  N. Weaver, *Quantum relations*, AMS, 2012.